

Multi-Head Monitoring of Metric Temporal Logic

Martin Raszyk, David Basin, Srđan Krstić, and Dmitriy Traytel

ETH zürich

Organisation of the Talk

Introduction — Monitoring Problem and Metric Temporal Logic

Online vs Multi-Head Monitoring

MTL Multi-Head Monitor

Evaluation

Future Work

Roadmap

Introduction — Monitoring Problem and Metric Temporal Logic

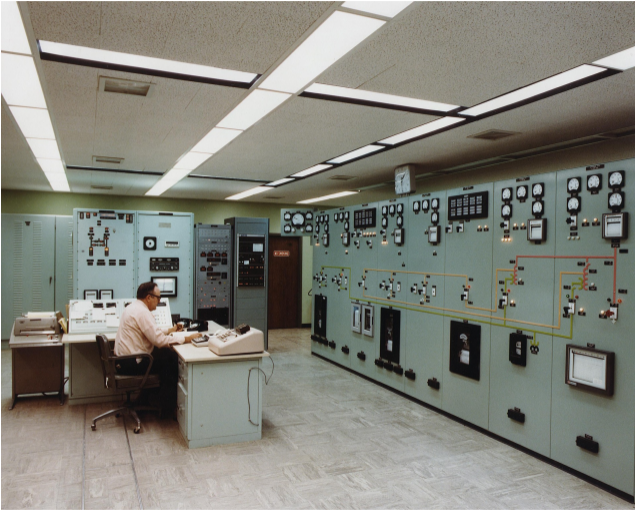
Online vs Multi-Head Monitoring

MTL Multi-Head Monitor

Evaluation

Future Work

Monitoring Problem



Example Property¹

Every alarm is followed by a shut down event in 10 time units unless all clear is sounded first.

¹Joël Ouaknine and James Worrell, FORMATS 2008

Linear Temporal Logic (with Past and Future)

Syntax:

$$\varphi = p \mid \neg\varphi \mid \varphi \vee \varphi \mid \bullet \varphi \mid \circ \varphi \mid \varphi S \varphi \mid \varphi U \varphi$$

p : atomic proposition

Linear Temporal Logic (with Past and Future)

Syntax:

$$\varphi = p \mid \neg\varphi \mid \varphi \vee \varphi \mid \bullet \varphi \mid \circ \varphi \mid \varphi S \varphi \mid \varphi U \varphi$$

p : atomic proposition

Semantics:



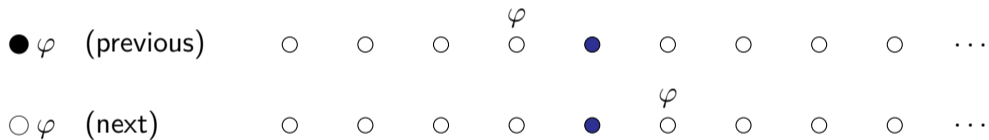
Linear Temporal Logic (with Past and Future)

Syntax:

$$\varphi = p \mid \neg\varphi \mid \varphi \vee \varphi \mid \bullet \varphi \mid \circ \varphi \mid \varphi S \varphi \mid \varphi U \varphi$$

p : atomic proposition

Semantics:



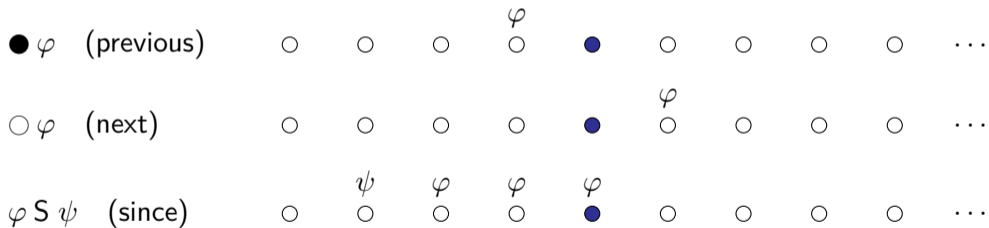
Linear Temporal Logic (with Past and Future)

Syntax:

$$\varphi = p \mid \neg\varphi \mid \varphi \vee \varphi \mid \bullet\varphi \mid \circ\varphi \mid \varphi S \varphi \mid \varphi U \varphi$$

p : atomic proposition

Semantics:



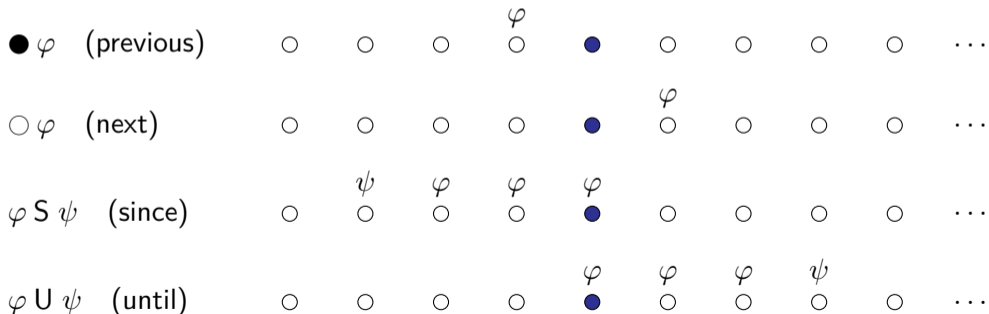
Linear Temporal Logic (with Past and Future)

Syntax:

$$\varphi = p \mid \neg\varphi \mid \varphi \vee \varphi \mid \bullet \varphi \mid \circ \varphi \mid \varphi S \varphi \mid \varphi U \varphi$$

p : atomic proposition

Semantics:



Example Property

Every alarm is followed by a shut down event in 10 time units unless all clear is sounded first.

Atomic propositions:

$$P = \{\text{alarm, shut_down, all_clear}\}$$

Example Property

*Every **alarm** is followed by a shut down event in 10 time units unless all clear is sounded first.*

Atomic propositions:

$$P = \{\text{alarm}, \text{shut_down}, \text{all_clear}\}$$

Example Property

*Every alarm is followed by a **shut down** event in 10 time units unless all clear is sounded first.*

Atomic propositions:

$$P = \{\text{alarm, shut_down, all_clear}\}$$

Example Property

*Every alarm is followed by a shut down event in 10 time units unless **all clear** is sounded first.*

Atomic propositions:

$$P = \{\text{alarm, shut_down, all_clear}\}$$

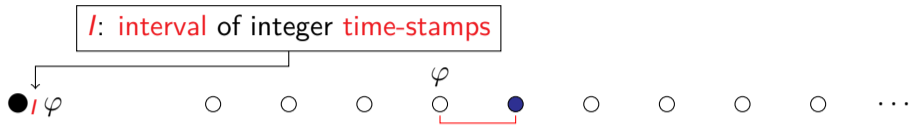
Example Property

Every alarm is followed by a shut down event in 10 time units unless all clear is sounded first.

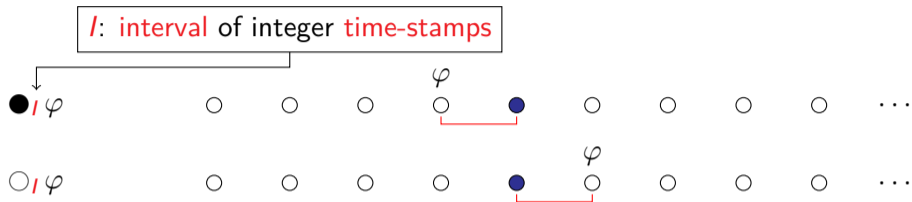
Atomic propositions:

$$P = \{\text{alarm, shut_down, all_clear}\}$$

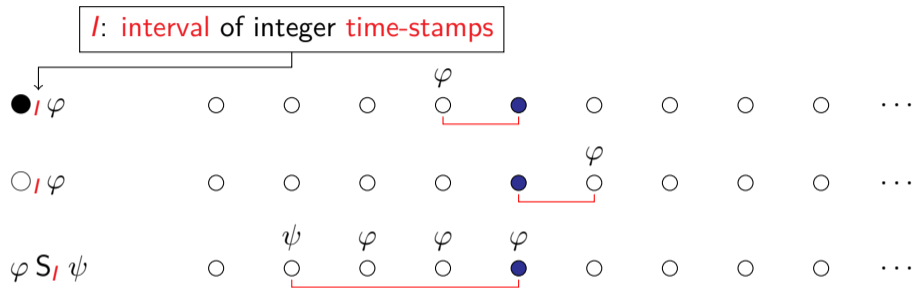
Metric Temporal Logic



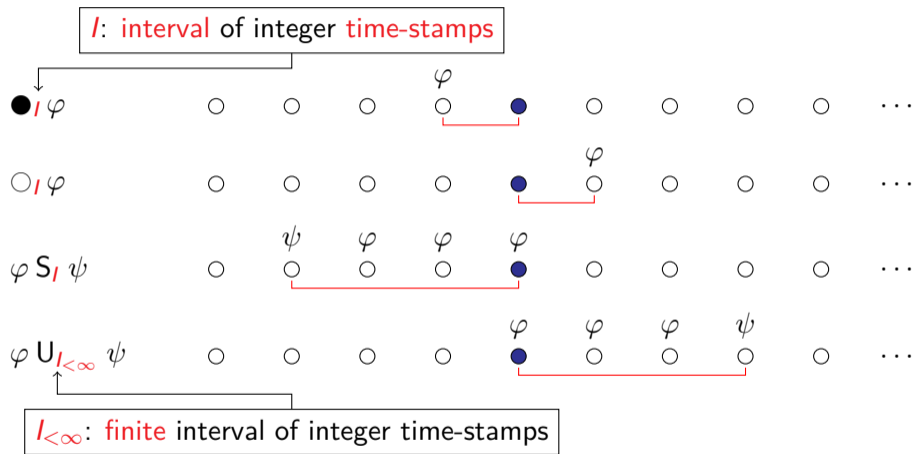
Metric Temporal Logic



Metric Temporal Logic



Metric Temporal Logic



Example Property

Every alarm is followed by a shut down event in 10 time units unless all clear is sounded first.

MTL formula:

$$\text{alarm} \rightarrow (\diamond_{[0,10]} \text{all_clear} \vee \diamond_{[10,10]} \text{shut_down})$$



MTL Monitoring Problem: Definition

Input MTL formula φ , time-stamped stream of events ρ .

Output *Stream* of Boolean values denoting whether $(\rho, i) \models \varphi$.

MTL Monitoring Problem: Definition

Input MTL formula φ , time-stamped stream of events ρ .

Output *Stream* of Boolean values denoting whether $(\rho, i) \models \varphi$.

Example:

$$\text{alarm} \rightarrow (\diamond_{[0,10]} \text{all_clear} \vee \diamond_{[10,10]} \text{shut_down})$$

Time-stamped stream ρ

Boolean output stream

MTL Monitoring Problem: Definition

Input MTL formula φ , time-stamped stream of events ρ .

Output *Stream* of Boolean values denoting whether $(\rho, i) \models \varphi$.

Example:

$$\text{alarm} \rightarrow (\diamond_{[0,10]} \text{all_clear} \vee \diamond_{[10,10]} \text{shut_down})$$

Time-stamped stream ρ

@0 all_clear

Boolean output stream

MTL Monitoring Problem: Definition

Input MTL formula φ , time-stamped stream of events ρ .

Output *Stream* of Boolean values denoting whether $(\rho, i) \models \varphi$.

Example:

$$\text{alarm} \rightarrow (\diamond_{[0,10]} \text{all_clear} \vee \diamond_{[10,10]} \text{shut_down})$$

Time-stamped stream ρ

@0 all_clear

Boolean output stream

@0 true

MTL Monitoring Problem: Definition

Input MTL formula φ , time-stamped stream of events ρ .

Output *Stream* of Boolean values denoting whether $(\rho, i) \models \varphi$.

Example:

$$\text{alarm} \rightarrow (\diamond_{[0,10]} \text{all_clear} \vee \diamond_{[10,10]} \text{shut_down})$$

Time-stamped stream ρ

@0 all_clear

@10 alarm

Boolean output stream

@0 true

MTL Monitoring Problem: Definition

Input MTL formula φ , time-stamped stream of events ρ .

Output Stream of Boolean values denoting whether $(\rho, i) \models \varphi$.

Example:

$$\text{alarm} \rightarrow (\diamond_{[0,10]} \text{all_clear} \vee \diamond_{[10,10]} \text{shut_down})$$

Time-stamped stream ρ

@0 all_clear

@10 alarm

@20 all_clear

Boolean output stream

@0 true

MTL Monitoring Problem: Definition

Input MTL formula φ , time-stamped stream of events ρ .

Output *Stream* of Boolean values denoting whether $(\rho, i) \models \varphi$.

Example:

$$\text{alarm} \rightarrow (\diamond_{[0,10]} \text{all_clear} \vee \diamond_{[10,10]} \text{shut_down})$$

Time-stamped stream ρ

@0 all_clear

@10 alarm

@20 all_clear

Boolean output stream

@0 true

@10 true

@20 true

MTL Monitoring Problem: Definition

Input MTL formula φ , time-stamped stream of events ρ .

Output Stream of Boolean values denoting whether $(\rho, i) \models \varphi$.

Example:

$$\text{alarm} \rightarrow (\diamond_{[0,10]} \text{all_clear} \vee \diamond_{[10,10]} \text{shut_down})$$

Time-stamped stream ρ

@0 all_clear

@10 alarm

@20 all_clear

@30 alarm

Boolean output stream

@0 true

@10 true

@20 true

MTL Monitoring Problem: Definition

Input MTL formula φ , time-stamped stream of events ρ .

Output Stream of Boolean values denoting whether $(\rho, i) \models \varphi$.

Example:

$$\text{alarm} \rightarrow (\diamond_{[0,10]} \text{all_clear} \vee \diamond_{[10,10]} \text{shut_down})$$

Time-stamped stream ρ

@0 all_clear

@10 alarm

@20 all_clear

@30 alarm

@45 all_clear

Boolean output stream

@0 true

@10 true

@20 true

MTL Monitoring Problem: Definition

Input MTL formula φ , time-stamped stream of events ρ .

Output Stream of Boolean values denoting whether $(\rho, i) \models \varphi$.

Example:

$$\text{alarm} \rightarrow (\diamond_{[0,10]} \text{all_clear} \vee \diamond_{[10,10]} \text{shut_down})$$

Time-stamped stream ρ

@0 all_clear

@10 alarm

@20 all_clear

@30 alarm

@45 all_clear

Boolean output stream

@0 true

@10 true

@20 true

@30 false

@45 true

Roadmap

Introduction — Monitoring Problem and Metric Temporal Logic

Online vs Multi-Head Monitoring

MTL Multi-Head Monitor

Evaluation

Future Work

Online Monitoring

Online monitor — algorithm reading event stream *once*, one event at a time.

Online Monitoring

Online monitor — algorithm reading event stream *once*, one event at a time.

Example:

$$\varphi = a \rightarrow \diamond_{[0,60]} b$$

Input stream ρ

Boolean output stream

Online Monitoring

Online monitor — algorithm reading event stream *once*, one event at a time.

Example:

$$\varphi = a \rightarrow \diamond_{[0,60]} b$$

Input stream ρ

@0 a

Boolean output stream

Online Monitoring

Online monitor — algorithm reading event stream *once*, one event at a time.

Example:

$$\varphi = a \rightarrow \diamond_{[0,60]} b$$

Input stream ρ

@0 a

Buffered events

@0 a

Boolean output stream

Online Monitoring

Online monitor — algorithm reading event stream *once*, one event at a time.

Example:

$$\varphi = a \rightarrow \diamond_{[0,60]} b$$

Input stream ρ

@0 a

@0

Buffered events

@0 a

Boolean output stream

Online Monitoring

Online monitor — algorithm reading event stream *once*, one event at a time.

Example:

$$\varphi = a \rightarrow \diamond_{[0,60]} b$$

Input stream ρ

@0 a

@0

Buffered events

@0 a

@0

Boolean output stream

Online Monitoring

Online monitor — algorithm reading event stream *once*, one event at a time.

Example:

$$\varphi = a \rightarrow \diamond_{[0,60]} b$$

Input stream ρ

@0 a

@0

@0 a

Buffered events

@0 a

@0

Boolean output stream

Online Monitoring

Online monitor — algorithm reading event stream *once*, one event at a time.

Example:

$$\varphi = a \rightarrow \diamond_{[0,60]} b$$

Input stream ρ

@0 a

@0

@0 a

Buffered events

@0 a

@0

@0 a

Boolean output stream

Online Monitoring

Online monitor — algorithm reading event stream *once*, one event at a time.

Example:

$$\varphi = a \rightarrow \diamond_{[0,60]} b$$

Input stream ρ

@0 a

@0

@0 a

@0

Buffered events

@0 a

@0

@0 a

Boolean output stream

Online Monitoring

Online monitor — algorithm reading event stream *once*, one event at a time.

Example:

$$\varphi = a \rightarrow \diamond_{[0,60]} b$$

Input stream ρ

@0 a

@0

@0 a

@0

Buffered events

@0 a

@0

@0 a

@0

Boolean output stream

Online Monitoring

Online monitor — algorithm reading event stream *once*, one event at a time.

Example:

$$\varphi = a \rightarrow \diamond_{[0,60]} b$$

Input stream ρ

@0 a

@0

@0 a

@0

@10 b

Buffered events

@0 a

@0

@0 a

@0

Boolean output stream

Online Monitoring

Online monitor — algorithm reading event stream *once*, one event at a time.

Example:

$$\varphi = a \rightarrow \diamond_{[0,60]} b$$

Input stream ρ

@0	a
@0	
@0	a
@0	
@10	b

Buffered events

Boolean output stream

@0	true
@0	true
@0	true
@0	true
@10	true

Online Monitoring

Online monitor — algorithm reading event stream *once*, one event at a time.

Example:

$$\varphi = a \rightarrow \diamond_{[0,60]} b$$

Input stream ρ

@0 a

@0

@0 a

@0

@80

Buffered events

@0 a

@0

@0 a

@0

Boolean output stream

Online Monitoring

Online monitor — algorithm reading event stream *once*, one event at a time.

Example:

$$\varphi = a \rightarrow \diamond_{[0,60]} b$$

Input stream ρ

@0 a

@0

@0 a

@0

@80

Buffered events

Boolean output stream

@0 false

@0 true

@0 false

@0 true

@80 true

Online Monitoring

Online monitor — algorithm reading event stream *once*, one event at a time.

Example:

$$\varphi = a \rightarrow \diamond_{[0,60]} b$$

Input stream ρ	Buffered events	Boolean output stream
@0 a		@0 false
@0		@0 true
@0 a		@0 false
@0		@0 true
@80		@80 true

Online monitor might need to **buffer** the **entire** trace in memory.

Related Work: AERIAL (TACAS 2017)

Idea

*Represent the Boolean output stream implicitly using **out-of-order** verdicts and **equivalence** verdicts.*

Related Work: AERIAL (TACAS 2017)

Idea

Represent the Boolean output stream implicitly using *out-of-order* verdicts and *equivalence* verdicts.

$$\varphi = a \rightarrow \diamond_{[0,60]} b$$

Input stream ρ

Buffered events

Output stream

Related Work: AERIAL (TACAS 2017)

Idea

Represent the Boolean output stream implicitly using *out-of-order* verdicts and *equivalence* verdicts.

$$\varphi = a \rightarrow \diamond_{[0,60]} b$$

Input stream ρ

@0 a

Buffered events

Output stream

Related Work: AERIAL (TACAS 2017)

Idea

Represent the Boolean output stream implicitly using *out-of-order* verdicts and *equivalence* verdicts.

$$\varphi = a \rightarrow \diamond_{[0,60]} b$$

Input stream ρ

@0 a

Buffered events

@0:0 a

Output stream

Related Work: AERIAL (TACAS 2017)

Idea

Represent the Boolean output stream implicitly using *out-of-order* verdicts and *equivalence* verdicts.

$$\varphi = a \rightarrow \diamond_{[0,60]} b$$

Input stream ρ

@0 a
@0

Buffered events

@0:0 a

Output stream

Related Work: AERIAL (TACAS 2017)

Idea

Represent the Boolean output stream implicitly using *out-of-order* verdicts and *equivalence* verdicts.

$$\varphi = a \rightarrow \diamond_{[0,60]} b$$

Input stream ρ

@0 a
@0

Buffered events

@0:0 a

Output stream

@0:1 true

Related Work: AERIAL (TACAS 2017)

Idea

Represent the Boolean output stream implicitly using *out-of-order* verdicts and *equivalence* verdicts.

$$\varphi = a \rightarrow \diamond_{[0,60]} b$$

Input stream ρ

@0 a

@0

@0 a

Buffered events

@0:0 a

Output stream

@0:1 true

Related Work: AERIAL (TACAS 2017)

Idea

Represent the Boolean output stream implicitly using *out-of-order* verdicts and *equivalence* verdicts.

$$\varphi = a \rightarrow \diamond_{[0,60]} b$$

Input stream ρ

@0 a

@0

@0 a

Buffered events

@0:0 a

Output stream

@0:1 true

@0:2 = @0:0

Related Work: AERIAL (TACAS 2017)

Idea

Represent the Boolean output stream implicitly using *out-of-order* verdicts and *equivalence* verdicts.

$$\varphi = a \rightarrow \diamond_{[0,60]} b$$

Input stream ρ

@0 a
@0
@0 a
@0

Buffered events

@0:0 a

Output stream

@0:1 true
@0:2 = @0:0

Related Work: AERIAL (TACAS 2017)

Idea

Represent the Boolean output stream implicitly using *out-of-order* verdicts and *equivalence* verdicts.

$$\varphi = a \rightarrow \diamond_{[0,60]} b$$

Input stream ρ

@0 a
@0
@0 a
@0

Buffered events

@0:0 a

Output stream

@0:1 true
@0:2 = @0:0
@0:3 true

Related Work: AERIAL (TACAS 2017)

Idea

Represent the Boolean output stream implicitly using *out-of-order* verdicts and *equivalence* verdicts.

$$\varphi = a \rightarrow \diamond_{[0,60]} b$$

Input stream ρ

@0 a
@0
@0 a
@0
@10 b

Buffered events

@0:0 a

Output stream

@0:1 true
@0:2 = @0:0
@0:3 true

Related Work: AERIAL (TACAS 2017)

Idea

Represent the Boolean output stream implicitly using *out-of-order* verdicts and *equivalence* verdicts.

$$\varphi = a \rightarrow \diamond_{[0,60]} b$$

Input stream ρ

@0	a
@0	
@0	a
@0	
@10	b

Buffered events

Output stream

@0:1	true
@0:2	= @0:0
@0:3	true
@0:0	true
@10:0	true

Related Work: AERIAL (TACAS 2017)

Idea

Represent the Boolean output stream implicitly using *out-of-order* verdicts and *equivalence* verdicts.

$$\varphi = a \rightarrow \diamond_{[0,60]} b$$

Input stream ρ

@0 a

@0

@0 a

@0

@80

Buffered events

@0:0 a

Output stream

@0:1 true

@0:2 = @0:0

@0:3 true

Related Work: AERIAL (TACAS 2017)

Idea

Represent the Boolean output stream implicitly using *out-of-order* verdicts and *equivalence* verdicts.

$$\varphi = a \rightarrow \diamond_{[0,60]} b$$

Input stream ρ

@0 a
@0
@0 a
@0
@80

Buffered events

Output stream

@0:1 true
@0:2 = @0:0
@0:3 true
@0:0 false
@80:0 true

Related Work: AERIAL (TACAS 2017)

Idea

Represent the Boolean output stream implicitly using *out-of-order* verdicts and *equivalence* verdicts.

$$\varphi = a \rightarrow \diamond_{[0,60]} b$$

Input stream ρ	Buffered events	Output stream
@0 a		@0:1 true
@0		@0:2 = @0:0
@0 a		@0:3 true
@0		@0:0 false
@80		@80:0 true

AERIAL only needs to **buffer** a **constant** number of events (independent of the trace).

Multi-Head Monitoring

Idea

Read multiple events from the event stream at once.

Multi-Head Monitoring

Idea

Read multiple events from the event stream at once.

$$\varphi = a \rightarrow \diamond_{[0,60]} b$$

Heads



Time-stamped stream ρ

@0 a
@0
@0 a
@0
@10 b

Boolean output stream

Multi-Head Monitoring

Idea

Read multiple events from the event stream at once.

$$\varphi = a \rightarrow \diamond_{[0,60]} b$$

Heads



Time-stamped stream ρ

@0 a

@0

@0 a

@0

@10 b

Boolean output stream

Multi-Head Monitoring

Idea

Read multiple events from the event stream at once.

$$\varphi = a \rightarrow \diamond_{[0,60]} b$$

Heads

Time-stamped stream ρ

Boolean output stream



@0 a

@0

@0 a

@0

@10 b

Multi-Head Monitoring

Idea

Read multiple events from the event stream at once.

$$\varphi = a \rightarrow \diamond_{[0,60]} b$$

Heads



Time-stamped stream ρ

@0 a

@0

@0 a

@0

@10 b

Boolean output stream

Multi-Head Monitoring

Idea

Read multiple events from the event stream at once.

$$\varphi = a \rightarrow \diamond_{[0,60]} b$$

Heads

Time-stamped stream ρ

Boolean output stream



@0 a

@0

@0 a

@0



@10 b

Multi-Head Monitoring

Idea

Read multiple events from the event stream at once.

$$\varphi = a \rightarrow \diamond_{[0,60]} b$$

Heads	Time-stamped stream ρ	Boolean output stream
	@0 a	@0 true
●	@0	
	@0 a	
	@0	
●	@10 b	

Multi-Head Monitoring

Idea

Read multiple events from the event stream at once.

$$\varphi = a \rightarrow \diamond_{[0,60]} b$$

Heads	Time-stamped stream ρ	Boolean output stream
	@0 a	@0 true
	@0	@0 true
●	@0 a	
	@0	
●	@10 b	

Multi-Head Monitoring

Idea

Read multiple events from the event stream at once.

$$\varphi = a \rightarrow \diamond_{[0,60]} b$$



Heads	Time-stamped stream ρ	Boolean output stream
	@0 a	@0 true
	@0	@0 true
	@0 a	@0 true
●	@0	
●	@10 b	

Multi-Head Monitoring

Idea

Read multiple events from the event stream at once.

$$\varphi = a \rightarrow \diamond_{[0,60]} b$$

Heads	Time-stamped stream ρ	Boolean output stream
	@0 a	@0 true
	@0	@0 true
	@0 a	@0 true
	@0	@0 true
 	@10 b	

Multi-Head Monitoring

Idea

Read multiple events from the event stream at once.

$$\varphi = a \rightarrow \diamond_{[0,60]} b$$

Heads	Time-stamped stream ρ	Boolean output stream
	@0 a	@0 true
	@0	@0 true
	@0 a	@0 true
	@0	@0 true
● ●	@10 b	@10 true

Multi-Head Monitoring

Idea

Read multiple events from the event stream at once.

$$\varphi = a \rightarrow \diamond_{[0,60]} b$$

Heads



Time-stamped stream ρ

@0 a

@0

@0 a

@0

@80 b

Boolean output stream

Multi-Head Monitoring

Idea

Read multiple events from the event stream at once.

$$\varphi = a \rightarrow \diamond_{[0,60]} b$$

Heads

Time-stamped stream ρ

Boolean output stream



@0 a

@0

@0 a

@0



@80 b

Multi-Head Monitoring

Idea

Read multiple events from the event stream at once.

$$\varphi = a \rightarrow \diamond_{[0,60]} b$$

Heads	Time-stamped stream ρ	Boolean output stream
	@0 a	@0 false
●	@0	
	@0 a	
	@0	
●	@80 b	

Multi-Head Monitoring

Idea

Read multiple events from the event stream at once.

$$\varphi = a \rightarrow \diamond_{[0,60]} b$$

Heads	Time-stamped stream ρ	Boolean output stream
	@0 a	@0 false
	@0	@0 true
●	@0 a	
	@0	
●	@80 b	

Multi-Head Monitoring

Idea

Read multiple events from the event stream at once.

$$\varphi = a \rightarrow \diamond_{[0,60]} b$$

Heads	Time-stamped stream ρ	Boolean output stream
	@0 a	@0 false
	@0	@0 true
	@0 a	@0 false
●	@0	
●	@80 b	

Multi-Head Monitoring

Idea

Read multiple events from the event stream at once.

$$\varphi = a \rightarrow \diamond_{[0,60]} b$$

Heads	Time-stamped stream ρ	Boolean output stream
	@0 a	@0 false
	@0	@0 true
	@0 a	@0 false
	@0	@0 true
● ●	@80 b	

Multi-Head Monitoring

Idea

Read multiple events from the event stream at once.

$$\varphi = a \rightarrow \diamond_{[0,60]} b$$

Heads	Time-stamped stream ρ	Boolean output stream
	@0 a	@0 false
	@0	@0 true
	@0 a	@0 false
	@0	@0 true
● ●	@80 b	@80 true

Multi-Head Monitoring

Idea

Read multiple events from the event stream at once.

$$\varphi = a \rightarrow \diamond_{[0,60]} b$$

Heads	Time-stamped stream ρ	Boolean output stream
	@0 a	@0 false
	@0	@0 true
	@0 a	@0 false
	@0	@0 true
● ●	@80 b	@80 true

Can produce **in-order explicit** Boolean verdicts in **constant** working memory (using one head for each atomic proposition).

Roadmap

Introduction — Monitoring Problem and Metric Temporal Logic

Online vs Multi-Head Monitoring

MTL Multi-Head Monitor

Evaluation

Future Work

Multi-Head Monitor's Structure

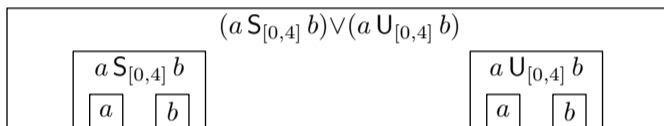
Example MTL formula:

$$(a S_{[0,4]} b) \vee (a U_{[0,4]} b)$$

Multi-Head Monitor's Structure

Example MTL formula:

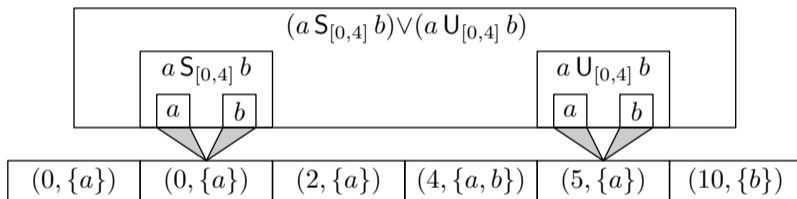
$$(a S_{[0,4]} b) \vee (a U_{[0,4]} b)$$



Multi-Head Monitor's Structure

Example MTL formula:

$$(a S_{[0,4]} b) \vee (a U_{[0,4]} b)$$



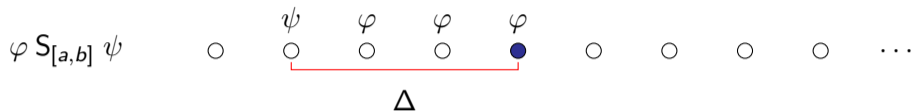
Since Operator Evaluation

Semantics:



Since Operator Evaluation

Semantics:

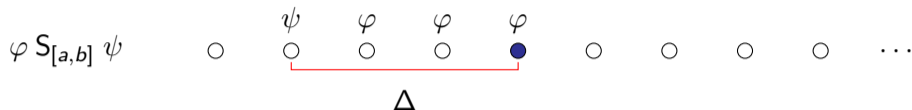


Idea (assuming $b \neq \infty$)

Store all the time-stamp differences $\Delta \in [0, b]$.

Since Operator Evaluation

Semantics:



Idea (assuming $b \neq \infty$)

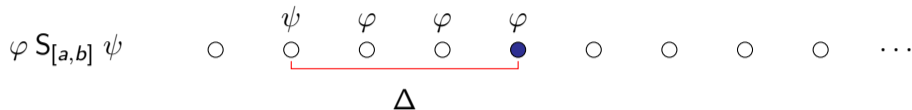
Store all the time-stamp differences $\Delta \in [0, b]$.

Need to store a set of natural numbers and (efficiently)

- ▶ query the maximum element,

Since Operator Evaluation

Semantics:



Idea (assuming $b \neq \infty$)

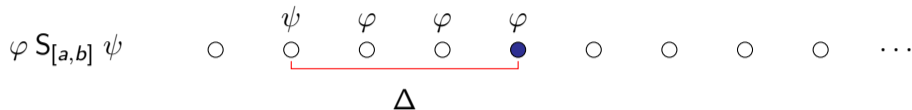
Store all the time-stamp differences $\Delta \in [0, b]$.

Need to store a set of natural numbers and (efficiently)

- ▶ query the maximum element,
- ▶ remove the maximum element,

Since Operator Evaluation

Semantics:



Idea (assuming $b \neq \infty$)

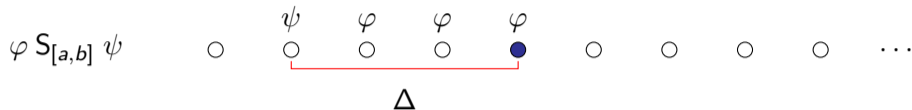
Store all the time-stamp differences $\Delta \in [0, b]$.

Need to store a set of natural numbers and (efficiently)

- ▶ query the maximum element,
- ▶ remove the maximum element,
- ▶ insert a zero,

Since Operator Evaluation

Semantics:



Idea (assuming $b \neq \infty$)

Store all the time-stamp differences $\Delta \in [0, b]$.

Need to store a set of natural numbers and (efficiently)

- ▶ query the maximum element,
- ▶ remove the maximum element,
- ▶ insert a zero,
- ▶ increase all elements by δ .

Delta Queue

Idea

Just store all inserted zeros and increases in a list and keep track of its sum!

Example:

Delta Queue

Idea

Just store all inserted zeros and increases in a list and keep track of its sum!

Example:

$$[] \quad (\Sigma = 0)$$

represents the multiset $\{\}$.

Delta Queue

Idea

Just store all inserted zeros and increases in a list and keep track of its sum!

Example:

$$[0] \quad (\Sigma = 0)$$

represents the multiset $\{0\}$.

Delta Queue

Idea

Just store all inserted zeros and increases in a list and keep track of its sum!

Example:

$$[0, 4] \quad (\Sigma = 4)$$

represents the multiset $\{4\}$.

Delta Queue

Idea

Just store all inserted zeros and increases in a list and keep track of its sum!

Example:

$$[0, 4, 2] \quad (\Sigma = 6)$$

represents the multiset $\{6\}$.

Delta Queue

Idea

Just store all inserted zeros and increases in a list and keep track of its sum!

Example:

$$[0, 4, 2, 0] \quad (\Sigma = 6)$$

represents the multiset $\{6, 0\}$.

Delta Queue

Idea

Just store all inserted zeros and increases in a list and keep track of its sum!

Example:

$$[0, 4, 2, 0^2] \quad (\Sigma = 6)$$

represents the multiset $\{6, 0^2\}$.

Delta Queue

Idea

Just store all inserted zeros and increases in a list and keep track of its sum!

Example:

$$[0, 4, 2, 0^2, 1] \quad (\Sigma = 7)$$

represents the multiset $\{7, 1^2\}$.

Delta Queue

Idea

Just store all inserted zeros and increases in a list and keep track of its sum!

Example:

$$[0, 4, 2, 0^2, 1, 0] \quad (\Sigma = 7)$$

represents the multiset $\{7, 1^2, 0\}$.

Delta Queue

Idea

Just store all inserted zeros and increases in a list and keep track of its sum!

Example:

$$[0, 4, 2, 0^2, 1, 0, 3] \quad (\Sigma = 10)$$

represents the multiset $\{10, 4^2, 3\}$.

Delta Queue

Idea

Just store all inserted zeros and increases in a list and keep track of its sum!

Example:

$$[0, 4, 2, 0^2, 1, 0, 3] \quad (\Sigma = 4)$$

represents the multiset $\{10, 4^2, 3\}$.

Delta Queue

Idea

Just store all inserted zeros and increases in a list and keep track of its sum!

Example:

$$[0^2, 1, 0, 3] \quad (\Sigma = 4)$$

represents the multiset $\{4^2, 3\}$.

Delta Queue

Idea

Just store all inserted zeros and increases in a list and keep track of its sum!

Example:

$$[0^2, 1, 0, 3] \quad (\Sigma = 4)$$

represents the multiset $\{4^2, 3\}$.

Each operation takes **amortized constant** time.

Delta Queue

Idea

Just store all inserted zeros and increases in a list and keep track of its sum!

Example:

$$[0^2, 1, 0, 3] \quad (\Sigma = 4)$$

represents the multiset $\{4^2, 3\}$.

Each operation takes **amortized constant** time and space complexity is linear in Σ .

Until Operator Evaluation

Semantics:



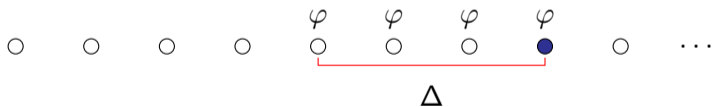
Until Operator Evaluation

Semantics:



Idea

Store all the time-stamp differences $\Delta \in [0, b]$ such that



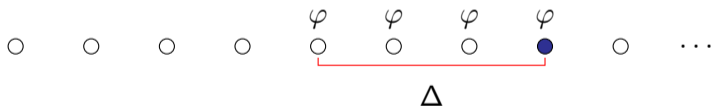
Until Operator Evaluation

Semantics:



Idea

Store all the time-stamp differences $\Delta \in [0, b]$ such that



We can reuse Delta Queue!

Space and Time Complexity

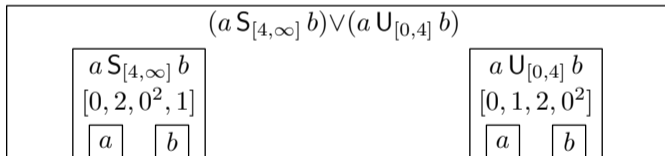
Size	$ \varphi $	number of MTL operators in φ	$(a S_{[4,\infty]} b) \vee (a U_{[0,4]} b)$
Temporal Size	$\ \varphi\ $	plus interval bounds	$(a S_{[4,\infty]} b) \vee (a U_{[0,4]} b)$

Space and Time Complexity

Size	$ \varphi $	number of MTL operators in φ	$(a S_{[4,\infty]} b) \vee (a U_{[0,4]} b)$
Temporal Size	$\ \varphi\ $	plus interval bounds	$(a S_{[4,\infty]} b) \vee (a U_{[0,4]} b)$

Theorem

Multi-head monitor's state for MTL formula φ requires $O(\|\varphi\|)$ registers storing time-stamps and indices into the trace.

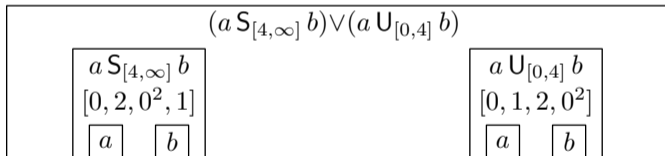


Space and Time Complexity

Size	$ \varphi $	number of MTL operators in φ	$(a S_{[4,\infty]} b) \vee (a U_{[0,4]} b)$
Temporal Size	$\ \varphi\ $	plus interval bounds	$(a S_{[4,\infty]} b) \vee (a U_{[0,4]} b)$

Theorem

Multi-head monitor's state for MTL formula φ requires $O(\|\varphi\|)$ registers storing time-stamps and indices into the trace.



Claim

Multi-head monitor for MTL formula φ runs in amortized time $O(|\varphi|)$ per event.

Roadmap

Introduction — Monitoring Problem and Metric Temporal Logic

Online vs Multi-Head Monitoring

MTL Multi-Head Monitor

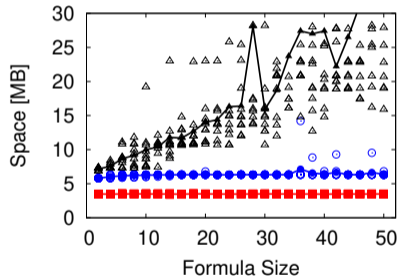
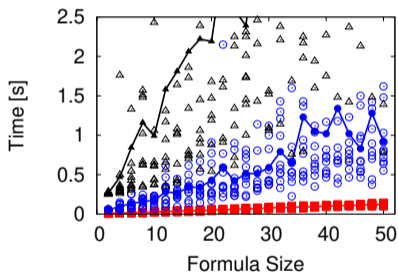
Evaluation

Future Work

Benchmarking Experiments — Average Case

Formula pseudo-random formula of given size.

Trace pseudo-random trace of fixed length.

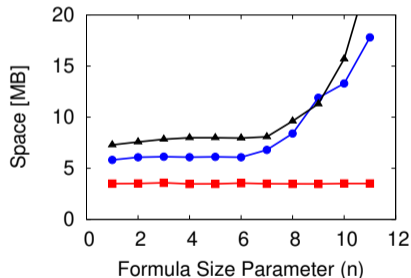
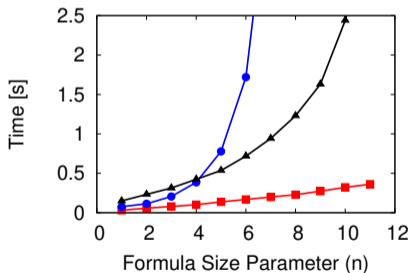


HYDRA —■— AERIAL —●— MONPOLY —▲—

Benchmarking Experiments — Worst Case

Formula MTL formula of size $O(n)$ encoding that a bit string $x \in \{0, 1\}^n$ precedes a distinguished labeled bit string $y \in \{0, 1\}^n$.

Trace all bit strings $x \in \{0, 1\}^n$ followed by a single bit string $y \in \{0, 1\}^n$ (pattern repeated to obtain a fixed length trace).



HYDRA —■— AERIAL —●— MONPOLY —▲—

Roadmap

Introduction — Monitoring Problem and Metric Temporal Logic

Online vs Multi-Head Monitoring

MTL Multi-Head Monitor

Evaluation

Future Work

Multi-Head Monitoring of Metric Dynamic Logic

Syntax:

$$\varphi = p \mid \neg\varphi \mid \varphi \vee \varphi \mid \langle r \rangle_I \mid \langle r \rangle_I \quad r = * \mid \varphi? \mid r + r \mid r \cdot r \mid r^*$$

²Martin Raszyk, David A. Basin, Dmitriy Traytel: From Nondeterministic to Multi-Head Deterministic Finite-State Transducers. ICALP 2019.

Multi-Head Monitoring of Metric Dynamic Logic

Syntax:

$$\varphi = p \mid \neg\varphi \mid \varphi \vee \varphi \mid \langle r \rangle \varphi \mid \langle r \rangle \varphi \quad r = * \mid \varphi? \mid r + r \mid r \cdot r \mid r^*$$

Building up on the following result²:

$$\begin{aligned} \mathcal{L}(2\text{DFT}) &= \mathcal{L}(f\text{-}2\text{NFT}) \\ &\quad \cup \\ \mathcal{L}(1\text{DFT}) &\subsetneq \mathcal{L}(f\text{-}1\text{NFT}) \subsetneq \mathcal{L}(\text{MH-}1\text{DFT}) \end{aligned}$$

²Martin Raszyk, David A. Basin, Dmitriy Traytel: From Nondeterministic to Multi-Head Deterministic Finite-State Transducers. ICALP 2019.